

On the pseudo-Frobenius numbers of numerical semigroups

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* $\mathbb{Z} = \{\dots, -2, -1, 0, 1, 2, \dots\}$

* $\mathbb{N} = \{0, 1, 2, \dots\}$

Definition

▶ S numerical semigroup = S cofinite submonoid of $(\mathbb{N}, +)$.

Definition

▶ $x \notin S$ is the Frobenius number of S if $x + s \in S$ for all $s \in \mathbb{N} \setminus \{0\}$.

▶ $x \notin S$ is a pseudo-Frobenius number of S if $x + s \in S$ for all $s \in S \setminus \{0\}$.

* $\text{PF}(S) = \{\text{Pseudo-Frobenius numbers of } S\}$

* $t(S) = \#\text{PF}(S) = \text{type of } S$

* $f(S) = \text{Frobenius number of } S$

* $f(S) = \max(\text{PF}(S)) = \max(\mathbb{Z} \setminus S)$

* $m(S) = \min(S \setminus \{0\})$

Motivation (1)

- * $f \in \mathbb{N} \setminus \{0\}$
- * $\mathcal{S}(f) = \{S \text{ numerical semigroup} \mid F(S) = f\}$
- * $\mathcal{S}(f) \neq \emptyset$ for all $f \in \mathbb{N} \setminus \{0\}$
- * There are algorithms to compute $\mathcal{S}(f)$

(Rosales, García-Sánchez, García-García, Jiménez-Madrid, J. Pure Appl. Algebra **189**, 2004)

- ** $\text{PF} = \{g_1, g_2, \dots, g_{n-1}, g_n\} \subseteq \mathbb{N} \setminus \{0\}$
- ** $\mathcal{S}(\text{PF}) = \{S \text{ numerical semigroup} \mid \text{PF}(S) = \text{PF}\}$
- ** $\mathcal{S}(\text{PF}) \neq \emptyset$ if $n = 1$
- ** $\mathcal{S}(\text{PF})$ may be empty if $n \geq 2$
- ** $\mathcal{S}(\text{PF})$ is finite when it is non-empty

Motivation (2)

Questions

1. Find conditions on the set PF that ensure that $\mathcal{S}(\text{PF}) \neq \emptyset$.
2. Find an algorithm to compute $\mathcal{S}(\text{PF})$.

Solved cases

- ▶ $\text{PF} = \{f\}$: symmetric numerical semigroups.
(Blanco, Rosales, Forum Math. **25**, 2013)
- ▶ $\text{PF} = \{f, \frac{f}{2}\}$: pseudo-symmetric numerical semigroups.
(Blanco, Rosales, Forum Math. **25**, 2013)
- ▶ Question 1 for $\text{PF} = \{g_1, g_2\}$.
(Robles-Pérez, Rosales, preprint)

Motivation (3)

Proposition

- ▶ Let g_1, g_2 be two positive integers such that $g_1 < g_2$. Let us take
- ▶ Let $g^* = \begin{cases} 2g_1 - g_2 & \text{if } g_2 \text{ is odd,} \\ g_1 - \frac{g_2}{2} & \text{if } g_2 \text{ is even.} \end{cases}$
- ▶ There exists a numerical semigroup S such that $\text{PF}(S) = \{g_1, g_2\}$ if and only if

$$\frac{g_2}{2} < g_1, \quad g^* \nmid g_1, \quad g^* \nmid g_2, \quad g^* \nmid (g_2 - g_1).$$

Example

- ▶ $\text{PF} = \{11, 15\}$ is an admissible set of pseudo-Frobenius numbers.
- ▶ $S = \langle 6, 7, 10 \rangle$ and $S = \langle 7, 9, 10, 12, 13 \rangle$.

Motivation (4)

Example

- ▶ $PF = \{9, 15\}$ is not a set of pseudo-Frobenius numbers.
($g^* = 2 \times 9 - 15 = 3$)
- ▶ $PF = \{9, 10, 15\}$ is a set of pseudo-Frobenius numbers.
 $S = \langle 4, 13, 14, 19 \rangle$.
($g_1^* = 2 \times 9 - 15 = 3$, $g_2^* = 2 \times 9 - 10 = 8$, $g_3^* = 2 \times 10 - 15 = 5$)

Example

- ▶ $PF = \{12, 18\}$ is not a set of pseudo-Frobenius numbers.
($g^* = 12 - \frac{18}{2} = 3$)
- ▶ $PF = \{12, 16, 18\}$ is a set of pseudo-Frobenius numbers.
 $S = \langle 7, 10, 13, 15, 19 \rangle$.
($g_1^* = 12 - \frac{18}{2} = 3$, $g_2^* = 12 - \frac{16}{2} = 4$, $g_3^* = 16 - \frac{18}{2} = 7$)

A simple idea to compute $\mathcal{S}(\text{PF})$: by filtering $\mathcal{S}(\max(\text{PF}))$

```
gap> pf := [17,27];;  
gap> nsf27 := NumericalSemigroupsWithFrobeniusNumber(27);;time;  
12062  
gap> Length(nsf27);  
16132  
gap> nspf1727 := Filtered(nsf27, s -> PseudoFrobeniusOfNumerical  
Semigroup(s) = pf);;time;Length(nspf1727);  
1313  
5
```

```
gap> pf := [19,29];;  
gap> nsf29 := NumericalSemigroupsWithFrobeniusNumber(29);;time;  
43281  
gap> Length(nsf29);  
34903  
gap> nspf1929 := Filtered(nsf29, s -> PseudoFrobeniusOfNumerical  
Semigroup(s) = pf);;time;Length(nspf1929);  
3171  
13
```

Alternatives

- * Determine the elements and the gaps (“forced” integers) of all numerical semigroups with the given set of pseudo-Frobenius numbers, obtaining a list of “free” integers. We then construct a tree in which branches correspond to assuming that free integers are either gaps or elements.
- ** Use irreducible numerical semigroups: from each irreducible numerical semigroup we start removing minimal generators with certain properties to build a tree whose leafs are the numerical semigroups we are looking for.

Forced and free integers

- * $\text{PF} = \{g_1, g_2, \dots, g_{n-1}, g_n\}$ a given set ($\subseteq \mathbb{N} \setminus \{0\}$)
- * $l \in \mathcal{GF}(\text{PF})$ if l is a gap of any $S \in \mathcal{S}(\text{PF})$
- * $l \in \mathcal{EF}(\text{PF})$ if $l \leq g_n + 1$ and it is an element of any $S \in \mathcal{S}(\text{PF})$

Proposition

- $\mathcal{S}(\text{PF}) \neq \emptyset$ if and only if $\mathcal{GF}(\text{PF}) \cap \mathcal{EF}(\text{PF}) = \emptyset$.

** $\text{Free}(\text{PF}) = \{1, \dots, g_n\} \setminus (\mathcal{GF}(\text{PF}) \cup \mathcal{EF}(\text{PF}))$

Example

- $\text{PF} = \{4, 17, 21\}$
- $\mathcal{GF}(\text{PF}) = \{1, 2, 3, 4, 5, 7, 9, 13, 17, 21\}$
- $\mathcal{EF}(\text{PF}) = \{8, 12, 14, 16, 18, 19, 20\}$
- $\text{Free}(\text{PF}) = \{6, 10, 11, 15\}$

Starting forced gaps (1)

Lemma

- ▶ Let S be a numerical semigroup.
- ▶ Then $m(S) \geq t(S) + 1$.

Lemma

- ▶ Let $\text{PF}(S) = \{g_1 < g_2 < \dots < g_{n-1} < g_n\}$, with $n > 1$.
- ▶ Let $i \in \{2, \dots, t(S)\}$ and $g \in \langle \text{PF}(S) \rangle$ with $g < g_i$.
- ▶ Then $g_i - g \in \text{gaps}(S) = \mathbb{N} \setminus S$.

* The set of positive divisors of

$$\text{PF}(S) \cup \{x \in \mathbb{N} \mid 1 \leq x \leq t(S)\} \cup$$

$$\{g_i - g \mid i \in \{2, \dots, t(S)\}, g \in \langle \text{PF}(S) \rangle, g_i > g\}$$

is a subset of $\text{gaps}(S)$.

Starting forced gaps (2)

- * $\text{PF} = \{g_1 < g_2 < \cdots < g_{n-1} < g_n\}$ ($n > 1$)
- * Let $\text{sfg}(\text{PF})$ be the set of positive divisors of $\text{PF} \cup \{x \in \mathbb{N} \mid 1 \leq x \leq n\} \cup \{g_i - g \mid i \in \{2, \dots, n\}, g \in \langle \text{PF} \rangle, g_i > g\}$.

Corollary

- ▶ *Let PF be a set of positive integers.*
- ▶ *Then $\text{sfg}(\text{PF}) \subseteq \mathcal{GF}(\text{PF})$.*

Example

- ▶ $\text{PF} = \{4, 9\} \Rightarrow \text{sfg}(\text{PF}) = \{1, 2, 3, 4, 5, 9\}$
- ▶ $\text{PF} = \{11, 17\} \Rightarrow \text{sfg}(\text{PF}) = \{1, 2, 3, 6, 11, 17\}$

Initial necessary condition

Lemma

- ▶ *Let S be a numerical semigroup.*
- ▶ *Then $x \in \mathbb{Z} \setminus S$ if and only if $f - x \in S$ for some $f \in \text{PF}(S)$.*

Lemma

- ▶ *Let $\text{PF}(S) = \text{PF} = \{g_1 < \dots < g_n\}$.*
- ▶ *Let $i \in \{2, \dots, n\}$ and $g \in \langle \text{PF} \rangle \setminus \{0\}$ with $g < g_i$.*
- ▶ *Then there exists $k \in \{1, \dots, n\}$ such that $g_k - (g_i - g) \in S$.*

Corollary

- ▶ $g_1 \geq g_n - g_{n-1}$.

Example

- ▶ $\text{PF} = \{4, 9\} \Rightarrow \mathcal{S}(\text{PF}) = \emptyset$

Forced elements: big forced elements

Lemma

- ▶ Let S be a numerical semigroup such that $m(S) = m$.
- ▶ Then $F(S) - i \in S \cup PF(S)$ for all $i \in \{1, \dots, m\}$.

Consequence

- ▶ Let $fg(PF) \subseteq \mathcal{GF}(PF)$.
- ▶ Let $m = \min(\mathbb{N} \setminus (fg(PF) \cup \{0\}))$.
- ▶ Then $\max(PF) - i \in \mathcal{EF}(PF) \cup PF$ for all $i \in \{1, \dots, m\}$.

Example

- ▶ $PF = \{4, 9\} \Rightarrow fg(PF) = sfg(PF) = \{1, 2, 3, 4, 5, 9\} \Rightarrow bfe(PF) = \{5, 6, 7, 8\}$
Contradiction!
- ▶ $PF = \{11, 17\} \Rightarrow fg(PF) = sfg(PF) = \{1, 2, 3, 6, 11, 17\} \Rightarrow$
 $bfe(PF) = \{14, 15, 16\}$

Forced elements: elements forced by exclusion (1)

Lemma

- ▶ Let S be a numerical semigroup.
- ▶ Then $x \in \mathbb{Z} \setminus S$ if and only if $f - x \in S$ for some $f \in \text{PF}(S)$.

Consequence

- ▶ Let $x \in \text{fg}(\text{PF}) \subseteq \mathcal{GF}(\text{PF})$.
- ▶ If there exists a unique $f^* \in \text{PF}(S)$ such that $f^* - x \notin \text{fg}(\text{PF})$, ...
- ▶ ... then $f^* - x \in \text{efe}_1(\text{PF}) \subseteq \mathcal{EF}(\text{PF})$.

Example

- ▶ $\text{PF} = \{11, 17\} \Rightarrow \text{fg}(\text{PF}) = \text{sfg}(\text{PF}) = \{1, 2, 3, 6, 11, 17\} \Rightarrow$
 $\text{efe}_1(\text{PF}) = \{5\}$

Forced elements: elements forced by exclusion (2)

Lemma

- ▶ Let S be a numerical semigroup.
- ▶ Then $x \in S$ if and only if $f - x \notin S$ for all $f \in \text{PF}(S)$.

Consequence

- ▶ Let $x \notin \text{fg}(\text{PF}) \cup \text{bfe}(\text{PF}) \cup \text{efe}_1(\text{PF})$.
- ▶ If $f - x \in \text{fg}(\text{PF})$, for every $f \in \text{PF}$, ...
- ▶ then $x \in \text{efe}_2(\text{PF}) \subseteq \mathcal{EF}(\text{PF})$.

Example

- ▶ $\text{PF} = \{11, 17\} \Rightarrow \text{fg}(\text{PF}) = \text{sfg}(\text{PF}) = \{1, 2, 3, 6, 11, 17\} \Rightarrow$
 $\text{bfe}(\text{PF}) \cup \text{efe}_1(\text{PF}) = \{5, 14, 15, 16\} \Rightarrow \text{efe}_2(\text{PF}) = \emptyset$
- * $\text{fe}(\text{PF}) = \langle \text{bfe}(\text{PF}) \cup \text{efe}_1(\text{PF}) \cup \text{efe}_2(\text{PF}) \rangle \cap [1, \max(\text{PF})] \subseteq \mathcal{EF}(\text{PF})$

Further forced gaps

Remark

- ▶ Let S be a numerical semigroup.
- ▶ Let $e \in S$ and $f \in \text{gaps}(S)$.
- ▶ Then $f - e < 0$ or $f - e \in \text{gaps}(S)$.

Consequence

- ▶ Let $\text{fg}(\text{PF}) \subseteq \mathcal{GF}(\text{PF})$ and $\text{fe}(\text{PF}) \subseteq \mathcal{EF}(\text{PF})$.
- ▶ Then $\Delta = (\text{fg}(\text{PF}) - e) \cap \mathbb{N} \subseteq \text{ffg}(\text{PF}) \subseteq \mathcal{GF}(\text{PF})$ for every $e \in \text{fe}(\text{PF})$.

Example

- ▶ $\text{PF} = \{11, 17\} \Rightarrow \text{sfg}(\text{PF}) = \{1, 2, 3, 6, 11, 17\} \Rightarrow$
 $\text{fe}(\text{PF}) = \{5, 10, 14, 15, 16\} \Rightarrow \Delta = \{7, 12\}$

* $\text{ffg}(\text{PF}) = \{\text{Positive divisors of } \text{sfg}(\text{PF}) \cup \Delta\} \subseteq \mathcal{GF}(\text{PF})$

Algorithm 1

- * PF
- * $PF \Rightarrow \text{sfg}(PF)$
- * $PF, \text{sfg}(PF) \Rightarrow \text{fe}(PF)$
- * $PF, \text{sfg}(PF), \text{fe}(PF) \Rightarrow \text{ffg}_1(PF)$
- * $PF, \text{ffg}_1(PF), \text{fe}(PF) \Rightarrow \text{ffe}_1(PF)$
- * $PF, \text{ffg}_1(PF), \text{ffe}_1(PF) \Rightarrow \text{ffg}_2(PF)$
- * $PF, \text{ffg}_2(PF), \text{ffe}_1(PF) \Rightarrow \text{ffe}_2(PF)$
- ⋮
- * Repeat until $\text{ffg}_{k+1}(PF) = \text{ffg}_k(PF)$ and $\text{ffe}_{k+1}(PF) = \text{ffe}_k(PF)$

- ** $\text{ffg}_k(PF) \subseteq \mathcal{GF}(PF), \text{ffe}_k(PF) \subseteq \mathcal{EF}(PF)$
- ** $[1, \max(PF)] \setminus (\text{ffg}_k(PF) \cup \text{ffe}_k(PF)) \supseteq \text{Free}(PF)$

Example

- * $PF = \{11, 17\}$
- * $PF = \{11, 17\} \Rightarrow \text{sfg}(PF) = \{1, 2, 3, 6, 11, 17\}$
- * $PF = \{11, 17\}, \text{sfg}(PF) = \{1, 2, 3, 6, 11, 17\} \Rightarrow \text{fe}(PF) = \{5, 10, 14, 15, 16\}$
- * $PF = \{11, 17\}, \text{sfg}(PF) = \{1, 2, 3, 6, 11, 17\} \text{fe}(PF) = \{5, 10, 14, 15, 16\} \Rightarrow$
 $\text{ffg}_1(PF) = \{1, 2, 3, 4, 6, 7, 11, 12, 17\}$
- * $PF = \{11, 17\}, \text{ffg}_1(PF) = \{1, 2, 3, 4, 6, 7, 11, 12, 17\}, \text{fe}(PF) =$
 $\{5, 10, 14, 15, 16\} \Rightarrow \text{ffe}_1(PF) = \{5, 10, 13, 14, 15, 16\}$
- * $PF = \{11, 17\}, \text{ffg}_1(PF) = \{1, 2, 3, 4, 6, 7, 11, 12, 17\}, \text{ffe}_1(PF) =$
 $\{5, 10, 13, 14, 15, 16\} \Rightarrow \text{ffg}_2(PF) = \{1, 2, 3, 4, 6, 7, 11, 12, 17\}$
- * $PF = \{11, 17\}, \text{ffg}_2(PF) = \{1, 2, 3, 4, 6, 7, 11, 12, 17\}, \text{ffe}_1(PF) =$
 $\{5, 10, 13, 14, 15, 16\} \Rightarrow \text{ffe}_2(PF) = \{5, 10, 13, 14, 15, 16\}$

- ** $\text{ffg}_2(PF) = \mathcal{GF}(PF) = \{1, 2, 3, 4, 6, 7, 11, 12, 17\}$
- ** $\text{ffe}_2(PF) = \mathcal{EF}(PF) = \{5, 10, 13, 14, 15, 16\}$
- ** $\text{Free}(PF) = \{8, 9\}$

Example

- * $PF = \{11, 17\}$, $\text{ffg}_2(PF) = \mathcal{GF}(PF) = \{1, 2, 3, 4, 6, 7, 11, 12, 17\}$
 $\text{ffe}_2(PF) = \mathcal{EF}(PF) = \{5, 10, 13, 14, 15, 16\}$, $\text{Free}(PF) = \{8, 9\}$

- ** Algorithm 1 with $\text{ffg}(PF) = \{1, 2, 3, 4, 6, 7, 11, 12, 17\}$ and
 $\text{ffe}(PF) = \{5, 8, 10, 13, 14, 15, 16\}$
 $S = \langle 5, 8, 14 \rangle = \{0, 5, 8, 10, 13, 14, 15, 16, 18, \rightarrow\}$

- ** Algorithm 1 with $\text{ffg}(PF) = \{1, 2, 3, 4, 6, 7, 8, 11, 12, 17\}$ and
 $\text{ffe}(PF) = \{5, 10, 13, 14, 15, 16\}$
 $S = \langle 5, 9, 13, 16 \rangle = \{0, 5, 9, 10, 13, 14, 15, 16, 18, \rightarrow\}$

Example

- * $PF = \{15, 20, 27, 35\}$
- * $sfg(PF) = \{1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 12, 15, 20, 27, 35\}$
- * $fe(PF) = \{19, 23, 25, 26, 28, 29, 30, 31, 32, 33, 34\}$
- * $ffg_1(PF) = \{1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 12, 15, 16, 20, 27, 35\}$
- * $ffe_1(PF) = \{19, 23, 25, 26, 28, 29, 30, 31, 32, 33, 34\}$
- * $ffg_2(PF) = \{1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 12, 15, 16, 20, 27, 35\}$
- * $ffe_2(PF) = \{19, 23, 25, 26, 28, 29, 30, 31, 32, 33, 34\}$

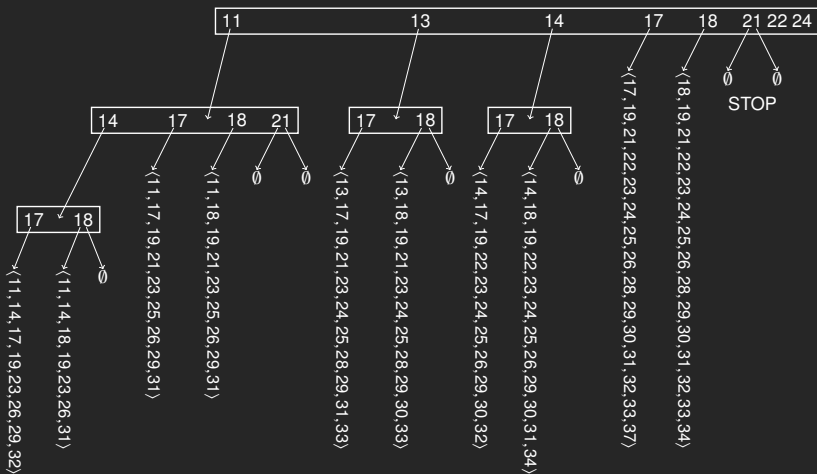
- ** $ffg_2(PF) = \{1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 12, 15, 20, 27, 35\} \subseteq \mathcal{GF}(PF)$
- ** $ffe_2(PF) = \{19, 23, 25, 26, 28, 29, 30, 31, 32, 33, 34\} \subseteq \mathcal{EF}(PF)$
- ** $Free(PF) \subseteq \{8, 9\}$

Example: $PF = \{15, 20, 27, 35\}$

$\{1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 12, 15, 16, 20, 27, 35\} \subseteq \mathcal{GF}(PF)$

$\{19, 23, 25, 26, 28, 29, 30, 31, 32, 33, 34\} \subseteq \mathcal{EF}(PF)$

$Free(PF) \subseteq \{11, 13, 14, 17, 18, 21, 22, 24\}$



Numerical semigroups with pseudo-Frobenius numbers $\{15, 20, 27, 35\}$

Running times

```
gap> pf:=[11,17];;  
gap> ans:=NumericalSemigroupsWithPseudoFrobeniusNumbers(pf);;  
gap> time;Length(ans);
```

0

2

```
gap> pf:=[17,27];;  
gap> ans:=NumericalSemigroupsWithPseudoFrobeniusNumbers(pf);;  
gap> time;Length(ans);
```

0

5

```
gap> pf:=[19,29];;  
gap> ans:=NumericalSemigroupsWithPseudoFrobeniusNumbers(pf);;  
gap> time;Length(ans);
```

16

13

Using irreducible numerical semigroups

Proposition

Let S be a nonirreducible numerical semigroup with $\text{PF}(S) = \{g_1 < \dots < g_k\}$, $k \geq 2$. Then there exists a chain

$$S = S_0 \subset S_1 = S \cup \{x_1\} \subset \dots \subset S_l = S \cup \{x_1, \dots, x_l\},$$

with

- (1) S_l irreducible,
- (2) $x_i = \max(S_l \setminus S_{i-1})$ for all $i \in \{1, \dots, l\}$,
- (3) $(g_k/2, g_k) \cap \text{PF}(S) \subset S_l$,
- (4) for every $i \in \{1, \dots, l\}$, x_i is a minimal generator of S_i such that $g_j - x_i \in S_i$ for some $j \in \{1, \dots, k\}$,
- (5) for every $i \in \{1, \dots, l\}$ and $f \in \text{PF}(S_i)$ with $f \neq g_k$, there exists $j \in \{1, \dots, k-1\}$ such that $g_j - f \in S_i$.

Example

* Irreducible numerical semigroups with $F(T_i) = 17$

** $T_1 = \langle 2, 19 \rangle$

** $T_2 = \langle 3, 10 \rangle$

** $T_3 = \langle 4, 6, 15 \rangle$

** $T_4 = \langle 4, 7 \rangle$

** $T_5 = \langle 4, 10, 11 \rangle$

** $T_6 = \langle 5, 8, 11, 14 \rangle$

** $T_7 = \langle 5, 9, 11, 13 \rangle$

** $T_8 = \langle 6, 7, 8 \rangle$

** $T_9 = \langle 6, 7, 9 \rangle$

** $T_{10} = \langle 6, 8, 10, 13, 15 \rangle$

** $T_{11} = \langle 6, 9, 10, 13, 14 \rangle$

** $T_{12} = \langle 7, 8, 11, 12, 13 \rangle$

** $T_{13} = \langle 7, 9, 11, 12, 13, 15 \rangle$

** $T_{14} = \langle 8, 10, 11, 12, 13, 14, 15 \rangle$

** $T_{15} = \langle 9, 10, 11, 12, 13, 14, 15, 16 \rangle$

* Numerical semigroups with $\text{PF}(S_i) = \{11, 17\}$

** $S_1 = T_6 \setminus \{11\} = \langle 5, 8, 14 \rangle$

** $S_2 = T_7 \setminus \{11\} = \langle 5, 9, 13, 16 \rangle$

THANK YOU VERY MUCH FOR YOUR ATTENTION!